

# Qudit Hypergraphs and Function Encoding

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with support from LVC Arnold research grants

# Why Quantum Computers?

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- Possibly solve new problems

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# Information Processing

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Further, **qudits** can be in a superposition of  $d$  different states:

$$|\psi\rangle = \alpha_0 |0\rangle + \alpha_1 |1\rangle + \alpha_2 |2\rangle + \dots + \alpha_{d-1} |d-1\rangle$$



The collective system of physically close states is described by the tensor-product of their states. A two qubit state would look something like:

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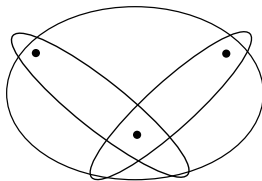
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One class of quantum states that has proven useful for quantum computation are hypergraph states

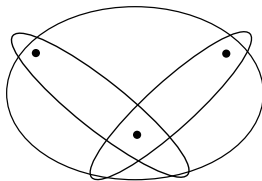
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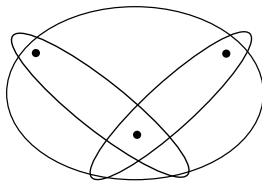
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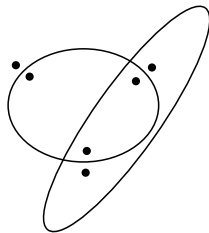
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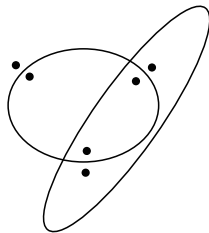
$$|\psi\rangle = |000\rangle + |001\rangle + |010\rangle - |011\rangle + |100\rangle + |101\rangle - |110\rangle - |111\rangle$$

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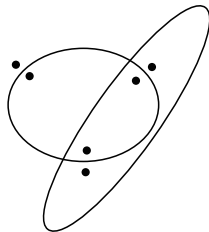
For a state with just three qudrits ( $d = 3$ ), there are 81 possible state vectors required to describe the state:

$$\begin{aligned} &= |000\rangle + |001\rangle + |002\rangle + |010\rangle + |011\rangle + |012\rangle + |020\rangle + |021\rangle + |022\rangle + |100\rangle + \\ &\omega |101\rangle + \omega |102\rangle + |110\rangle + \omega^2 |111\rangle + |112\rangle + |120\rangle + |121\rangle + \omega^2 |122\rangle + \\ &|200\rangle + \omega |201\rangle + \omega |202\rangle + |210\rangle + |211\rangle + \omega^2 |212\rangle + |220\rangle + \omega^2 |221\rangle + |222\rangle \end{aligned}$$



## Function Encoding (cont.)

Alternatively, these states can be described with uniquely encoded functions (mod  $d$ ).

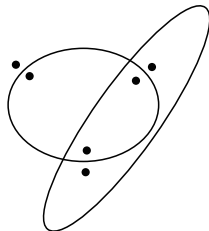


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The function,  $f$ , determines the phase for each individual element of the superposition. For example:

$$f(1, 2, 2) = 1 \cdot 2 \cdot 2 + 1^2 \cdot 2^2 = 2 \longrightarrow \omega^2 |122\rangle$$

# Encoding Process

The general map for encoding a hypergraph as a function is as follows:

$$\begin{array}{ccccc} \text{state} & & \text{c-vec} & & \text{a-vec} \\ |\psi\rangle & \longrightarrow & |c\rangle & \longrightarrow & |a\rangle \\ & \text{log}\omega & & S^{-1}\text{matrix} & \end{array}$$

Where the a vector ( $|a\rangle$ ) describes the coefficients of the state's function.  
Returning to the state is similar,

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The step from  $|c\rangle$  to  $|\psi\rangle$  is an entry-wise operation on the elements of a vector, while step  $|a\rangle$  to  $|c\rangle$  is achieved by multiplication of the  $S$  matrix

# Encoding Process (cont.)

For example:

$$|\psi\rangle = |00\rangle + |01\rangle + |02\rangle + |10\rangle + \omega |11\rangle + \omega^2 |12\rangle + |20\rangle + \omega^2 |21\rangle + \omega |22\rangle$$

↓

$$|\psi\rangle = [1 \quad 1 \quad 1 \quad 1 \quad \omega \quad \omega^2 \quad 1 \quad \omega^2 \quad \omega]$$

↓  $\log_\omega$

$$|c\rangle = [0 \quad 0 \quad 0 \quad 0 \quad 1 \quad 2 \quad 0 \quad 2 \quad 1]$$

↓  $S^{-1}$  Matrix

$$|a\rangle = [0 \quad 0 \quad 0 \quad 0 \quad 1 \quad 0 \quad 0 \quad 0 \quad 0]$$

$$x^0y^0 \mid x^0y^1 \mid x^0y^2 \mid x^1y^0 \mid x^1y^1 \mid x^1y^2 \mid x^2y^0 \mid x^2y^1 \mid x^2y^2$$

$$f = xy$$

# S Matrix

The  $S$  matrix has the form of a mod  $d$  Vandermonde matrix:

$$S_d = \begin{bmatrix} 0^0 & 0^1 & 0^2 & \dots & 0^{d-1} \\ 1^0 & 1^1 & 1^2 & \dots & 1^{d-1} \\ 2^0 & 2^1 & 2^2 & \dots & 2^{d-1} \\ 3^0 & 3^1 & 3^2 & \dots & 3^{d-1} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ (d-1)^0 & (d-1)^1 & (d-1)^2 & \dots & (d-1)^{d-1} \end{bmatrix} \pmod{d}$$

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For example, for  $d=5$

$$S_5 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 1 & 1 & 1 & 1 & 1 \\ 1 & 2 & 4 & 3 & 1 \\ 1 & 3 & 4 & 2 & 1 \\ 1 & 4 & 1 & 4 & 1 \end{bmatrix} \pmod{5}$$

# Generalized Local Pauli Matrices

The Paulis are a basis for  $d \times d$  operators (actions) on individual qudits.

They can be used to describe all local actions on a state. For qubits ( $d = 2$ ), they are standard:

$$\mathbf{I} = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} \quad X = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \quad Y = \begin{bmatrix} 0 & -i \\ i & 0 \end{bmatrix} \quad Z = \begin{bmatrix} 1 & 0 \\ 0 & -1 \end{bmatrix}$$



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We've considered one possible generalization of the Paulis for local  $d$  dimensional actions, the Ps (generalized X) and Qs (generalized Z).

For example:

$$P_{01} = \begin{bmatrix} 0 & 1 & 0 & 0 & 0 \\ 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \quad Q_0 = \begin{bmatrix} \omega & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

# Generalized Local Pauli Matrices (cont.)

The  $P$ s act as subspace permutations and the  $Q$ s act as local phase shifts on the state vector ( $|\psi\rangle$ ).

$$P_{24} = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 & 0 \\ 0 & 0 & 1 & 0 & 0 \end{bmatrix} \quad Q_3 = \begin{bmatrix} 1 & 0 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & \omega & 0 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix}$$

For any dimension  $d$ , there are  $d! - 1$   $P$ s and  $d$   $Q$ s.

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Pauli operators are important for error-correcting applications.

# GLP Actions on Encoded Functions

In order to see the effects of the generalized Paulis on encoded functions, we must consider the map of encoding,

$$|a'\rangle \leftarrow S^{-1} \leftarrow \log_{\omega} \leftarrow |\psi\rangle \leftarrow \exp_{\omega} \leftarrow S \leftarrow |a\rangle$$

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For example,

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$$f(x, y) \equiv 0, xy, xy^2, x^2y, xy^2 + x^2y, x^2y^2, 2x^2y^2, x^2y^2 + xy, 2x^2y^2 + xy$$

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All states  $n=2$ ,  $d=3$  can be transformed into one of the above functions with the GLP

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Stabilizers are important in application and theory.

## GLP Stabilizers (cont.)

One family of stabilizers found to work for any  $d$  or number of qudits is the 'Virginia Reel'

$$P_{VR} = \begin{bmatrix} 0 & 1 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \\ 1 & 0 & 0 & 0 & & 0 \end{bmatrix}$$



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And its transpose,

$$P_{VR}^T = \begin{bmatrix} 0 & 0 & 0 & 0 & \cdots & 1 \\ 1 & 0 & 0 & 0 & \cdots & 0 \\ 0 & 1 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \\ 0 & 0 & 0 & 0 & & 0 \end{bmatrix}$$

## GLP Stabilizers (cont.)

One family of stabilizers found to work for any  $d$  or number of qudits is the 'Virginia Reel'

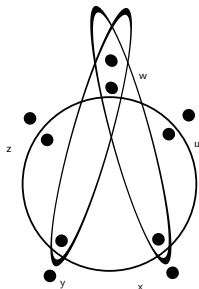
$$P_{VR} = \begin{bmatrix} 0 & 1 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 1 & 0 & \cdots & 0 \\ 0 & 0 & 0 & 1 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \\ 1 & 0 & 0 & 0 & & 0 \end{bmatrix}$$

And its transpose,

$$P_{VR}^T = \begin{bmatrix} 0 & 0 & 0 & 0 & \cdots & 1 \\ 1 & 0 & 0 & 0 & \cdots & 0 \\ 0 & 1 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \vdots & \ddots & \\ 0 & 0 & 0 & 0 & & 0 \end{bmatrix}$$

These will stabilize any function of the form  $f=(x+y)\cdot(\text{other variables})$

# GLP Stabilizers (cont.)



$$f(x, y, z, w, u) = (x + y)(zu + w^2)$$

The Virginia Reel is exponentiated from a hermitian matrix ( $M$ ) and a certain constant  $2\pi i/d$  ( $P_{VR} = e^{(2\pi i/d)M}$ ). Its transpose/inverse is exponentiated from the same hermitian matrix and the negative of the same constant ( $P_{VR}^T = e^{(-2\pi i/d)M}$ ). However, for any real  $t$ ,  $e^{itM} \otimes e^{-itM}$  is also a stabilizer for functions of the form  $f = (x+y) \cdot (\text{other variables})$ . Thus, there is an infinite family of stabilizers for such functions.

# Thank You!

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